

TWO STORIES THAT CAN BE TOLD IN PARALLEL.

I TAKUYAMA $q=0$ MODELS

PRACTIC MONOID; MULTIPLICATION OF CRYSTALS.

RSK = ROBINSON-SCHENSTED-KNUTH

BASIC ALGORITHM IN TABLEAU COMBINATORICS

BRACKETING: IN APPLYING CRYSTAL OPERATIONS WHAT DOES e_i, f_i DO?

II PIPEDREAMS

WEYL GROUP (FOCUS ON REDUCED ELEMENTS)

EDLMAN-GREENE INSERTION.

BRACKETING: IN APPLYING CRYSTAL OPERATIONS WHAT DOES e_i, f_i DO?

GOALS: BRACKETING IN CASE I.

INTRODUCE EDLMAN-GREENE AND USE IT TO PROVE STANLEY'S THM:

OF REDUCED WORDS = EXPLICIT NUMBER.
FOR $w_0 \in S_n$

$$\mathbb{B} = \boxed{1} \xrightarrow{1} \boxed{2} \xrightarrow{2} \boxed{3} \rightarrow \dots \xrightarrow{n-1} \boxed{n}$$

= STANDARD CRYSTAL FOR $GL(n)$

= COMBINATORIAL SUBSTITUTE FOR
STANDARD MODULE E^λ OF $GL(n)$.

$$\otimes^k \mathbb{B} = \bigoplus_{\lambda \vdash k} \mathbb{B}_\lambda$$

$\lambda \vdash k$

\uparrow

λ A PARTITION OF k ,

$$\text{IF } k=3 \quad f(\mathbb{B}) = f((1,1,1)) = 1$$

$$f((2,1)) = 2$$

IN GENERAL THERE IS AN IRR REP'N OF
 S_n DETERMINED $\prod_{\lambda} S_{\lambda}$ OF DIM'N f_{λ} .

$1, 1, 2$ = DEC. OF IRR REPS OF S_3

IN REP'N THEORY OF $GL(n)$

SCHUR-WEYL DUALITY:

$$\otimes^k \mathbb{C}^n = \mathbb{C}^n \otimes \dots \otimes \mathbb{C}^n$$

HAS COMMUTING REPS OF S_k AND $GL(n)$

$$\otimes^k \mathbb{C}^n = \bigoplus_{\substack{\lambda \vdash k \\ \ell(\lambda) \leq n}} \pi_{\lambda}^{S_k} \otimes \pi_{\lambda}^{GL(n)},$$

$n=3$ THIS SHOWS

$$\otimes^3 \mathbb{C}^n = \text{SYM}^3(\mathbb{C}^n) \oplus \wedge^3 \mathbb{C}^n \oplus 2 \pi_{(2,1,0,\dots)}^{GL(n)}$$

THIS EXPLAINS THE REASON THE REP'N THEORY OF S_k IS INVOLVED.

BY ANALOGY

$$\otimes^k \mathbb{B} = \bigoplus_{\substack{\lambda \vdash k \\ \ell(\lambda) \leq n}} p_{\lambda} \cdot \mathbb{B}_{\lambda}.$$

$$\mathbb{B}_\lambda \hookrightarrow \otimes^k \mathbb{B}$$

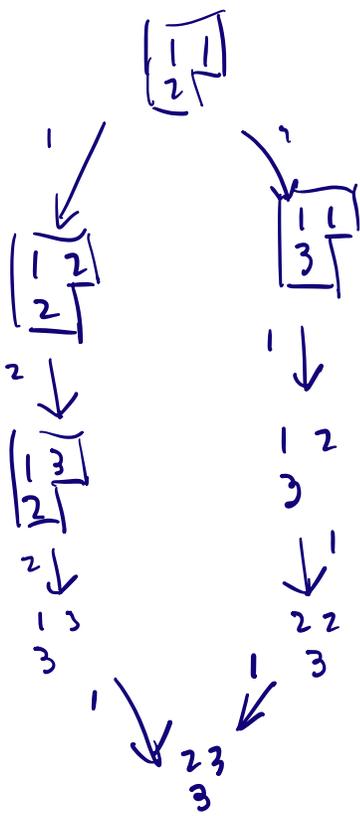
$$T \rightsquigarrow R_N \otimes R_{N-1} \otimes \dots \otimes R_1$$

R_1, \dots, R_N

ROWS.

$$R_i = \overline{[a_1 | a_2 | \dots | a_{\lambda_i}]} \hookrightarrow \overline{[a_1]} \otimes \overline{[a_2]} \otimes \dots$$

$$a_1 \leq a_2 \leq \dots \leq a_n.$$



$$\hookrightarrow \otimes^3 \mathbb{B}.$$

$$\overline{\begin{bmatrix} a & b \\ c \end{bmatrix}} \rightsquigarrow \overline{[c]} \otimes \overline{[a]} \otimes \overline{[b]}.$$

$$a \leq b$$

$$a < c.$$

THIS GIVES 8 ELEMENTS.

BRACKETING OPERATION DESCRIBES
 PROMOTION EXACTLY.

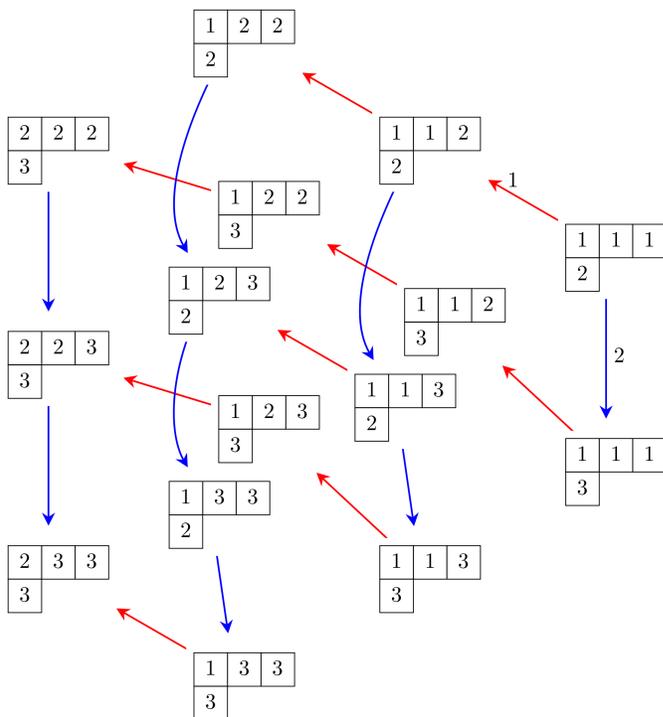
TO COMPUTE e_i OR f_i OF

$$T = \boxed{a_1} \otimes \boxed{a_2} \otimes \dots \otimes \boxed{a_n}$$

f_i CHANGES SOME $n \rightarrow n+1$

e_i " " $n+1 \rightarrow n$

EXCEPT SOMETIMES $f_i(T) = 0$



$GL(n)$

\downarrow

$U_q(\mathfrak{gl}(n))$

$$f_i = \begin{pmatrix} 0 & & \\ & 0 & \\ & & 0 \end{pmatrix}$$

$$e_i = \begin{pmatrix} & 1 & \\ & & \\ & & 0 \end{pmatrix}$$

$$V(\mu) = \{v \in V \mid H \cdot v = \lambda(H)v\}$$

$$V(\mu) \begin{array}{c} \xrightarrow{P_i} \\ \xleftarrow{e_i} \end{array} V(\mu - \alpha_i)$$

$$\boxed{a_1} \otimes \boxed{a_2} \otimes \dots \otimes \boxed{a_n}$$

f_i or e_i : IGNORE ALL ENTRIES $\neq i, i+1$

$$\boxed{3} \otimes \boxed{2} \otimes \boxed{2} \otimes \boxed{1} \otimes \boxed{1} \otimes \boxed{3} \otimes \boxed{1} \otimes \boxed{2}$$

\boxed{m} IS INVISIBLE IF $m \neq i, i+1$

$\boxed{i+1} \otimes \boxed{i}$ IS INVISIBLE.

$$f_i(\boxed{m} \otimes x) = \boxed{m} \otimes f_i(x)$$

$$f_i(\boxed{i+1} \otimes \boxed{i} \otimes x) = \boxed{i+1} \otimes \boxed{i} \otimes f_i(x)$$

$$\boxed{i+1} \otimes \boxed{m} \otimes \boxed{i}$$

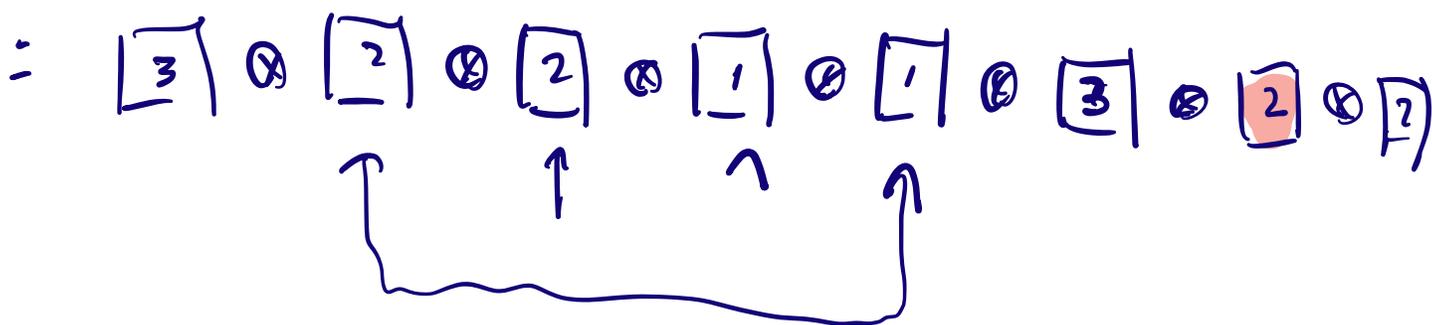
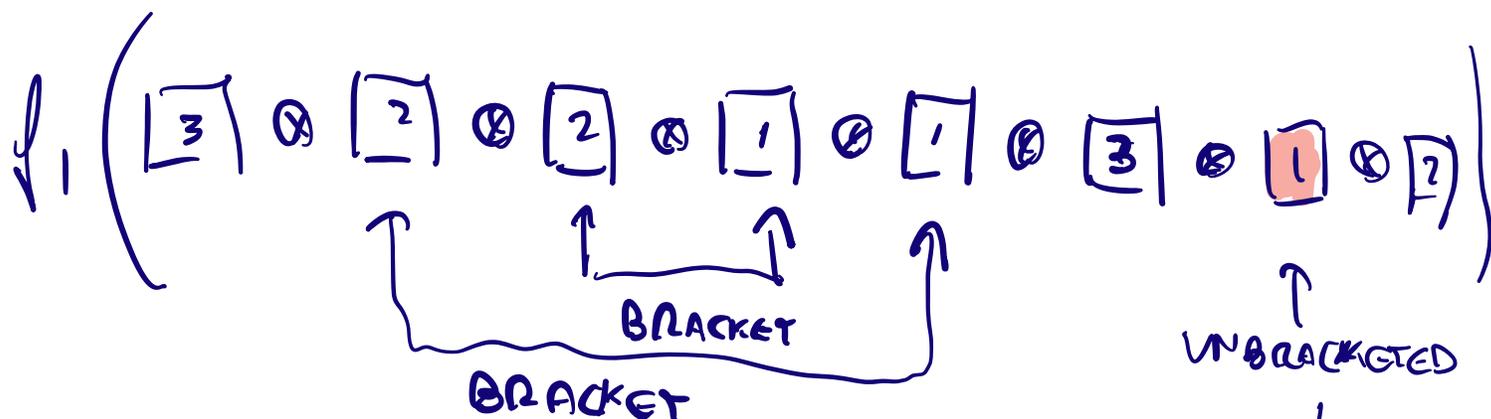
INVISIBLE
on $i, i+1$.

LOOK FOR $\boxed{i+1} \otimes (\text{INVISIBLE STUFF}) \otimes \boxed{i}$

↑ ↑

BRACKET.

$i = 1$



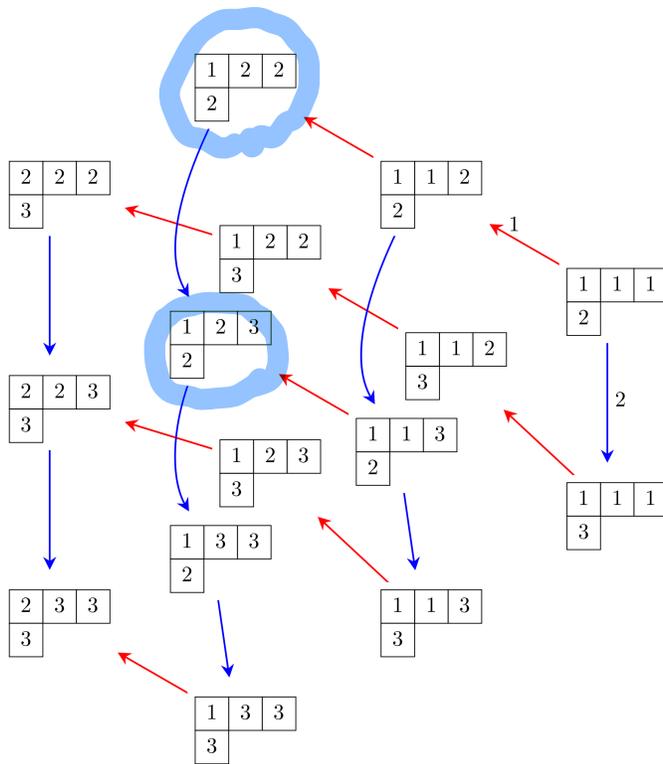
RULE: AFTER BRACKETING f_i CHANGES
RIGHT MOST UNBRACKETED $i \rightarrow i+1$.

IF ALL 'N'S' BRACKETED $f_i(T) = 0$.

IF ALL $i+1$ ARE BRACKETED $e_i(\tau) = 0$

OTHERWISE f_i SENDS LEF MOST

$$\boxed{i+1} \rightarrow \boxed{i}$$



$$f_2 \left(\begin{array}{ccc} 1 & 2 & 2 \\ 2 & & \end{array} \right) = f_2 \left(\begin{array}{cccc} 2 & 1 & 2 & 2 \\ \uparrow & \uparrow & & \end{array} \right)$$

$$2 \otimes 1 \otimes 2 \otimes 3 =$$

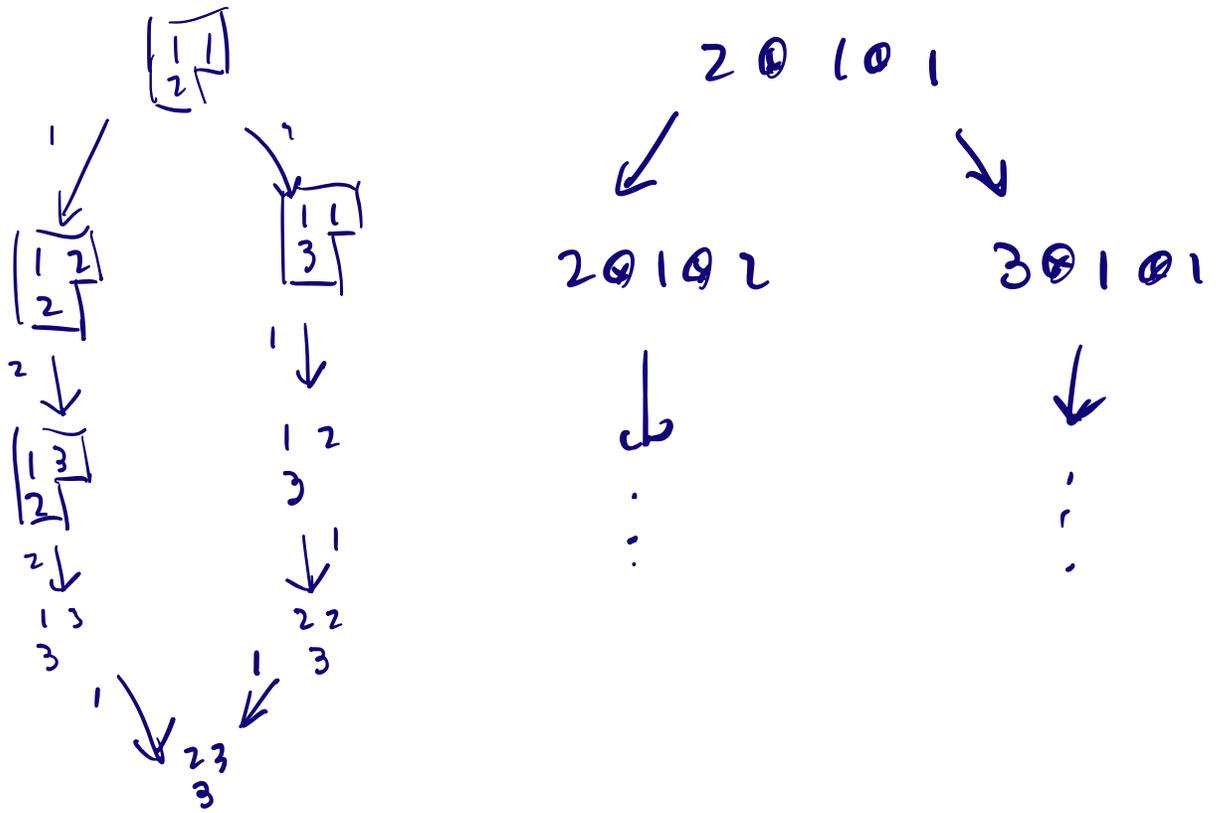
$$\left(\begin{array}{ccc} 1 & 2 & 3 \\ 2 & & \end{array} \right)$$

THE PLACTIC MONOID

(LASCoux & SCHUTZENBERGER -
 BASIC ALGORITHMS ARE DUE TO
 SCHENSTED AND KNUTH.)

DEFINE AN EQUIVALENCE RELATION

ON $\mathbb{Z}^n \times \mathbb{B} = \mathbb{A}$ IRREDUCIBLE CRYSTALS.



THERE MAY BE OTHER SUBCRYSTALS
OF $\otimes^k \mathbb{B}$ THAT ARE ISOMORPHIC

BUT NOT IN THE IMAGES OF
THE EMBEDDINGS $\mathbb{B}_\lambda \rightarrow \otimes^k \mathbb{B}$.

$$1 \otimes 2 \otimes 1$$

$$e_1 \left(\underbrace{1 \otimes 2 \otimes 1}_{\text{BRACKET}} \right) = 0 \quad \text{SINCE NO UNBRACKETED 2}$$

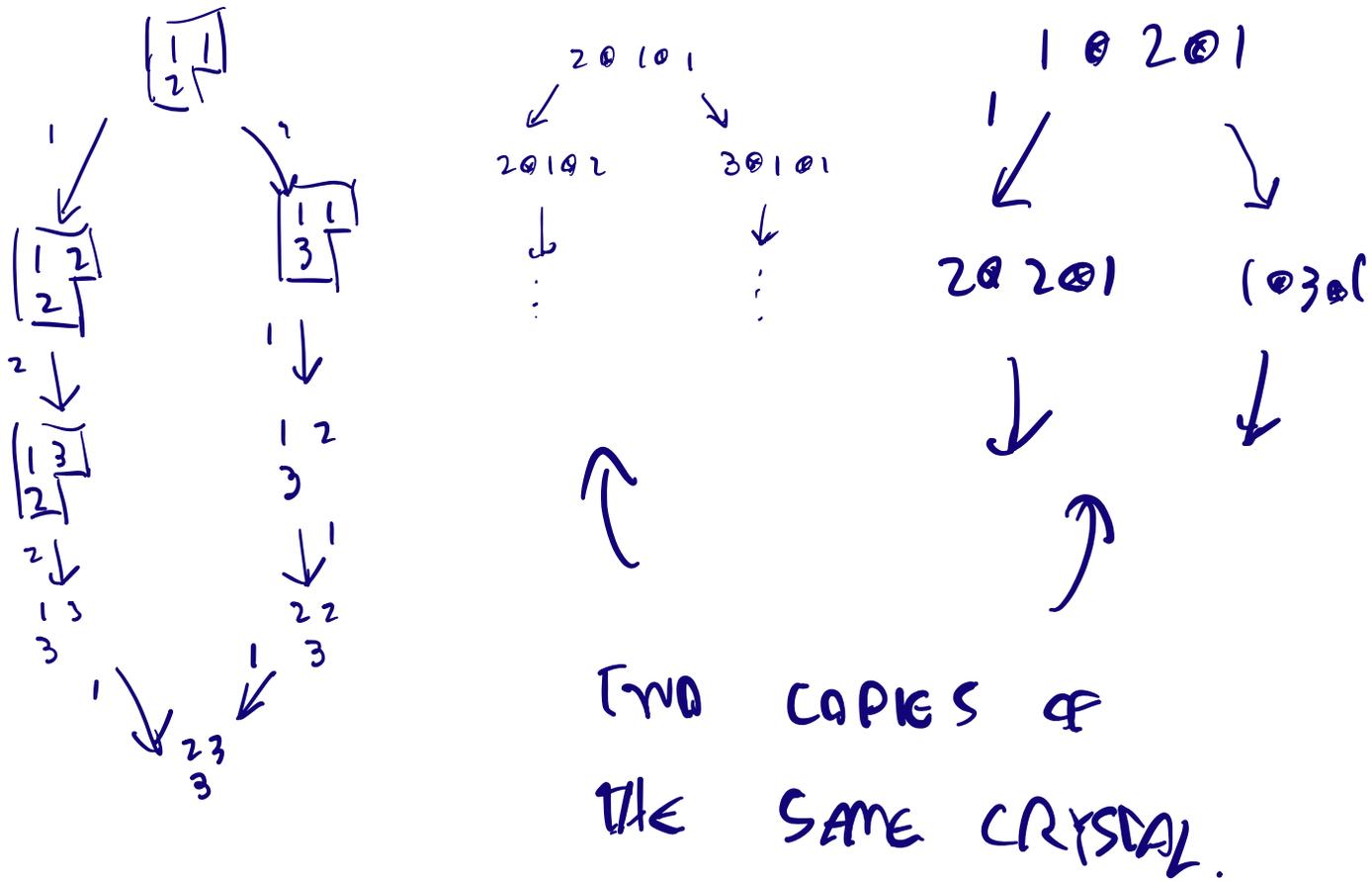
$$e_2 (1 \otimes 2 \otimes 1) = 0 \quad \text{NO UNBRACKETED 3.}$$

$$e_i (1 \otimes 2 \otimes 1) = 0 \quad \text{ALL } i$$

SO $1 \otimes 2 \otimes 1$ IS A HIGHEST WEIGHT ELEM.

IT GENERATES A CRYSTAL OF HIGHEST
WEIGHT $(2, 1, 0)$

I.E. A CRYSTAL ISOMORPHIC TO $\mathbb{B}_{(2,1,0)}$



Both embedded in $\mathbb{Q}^3 \mathbb{B}$.

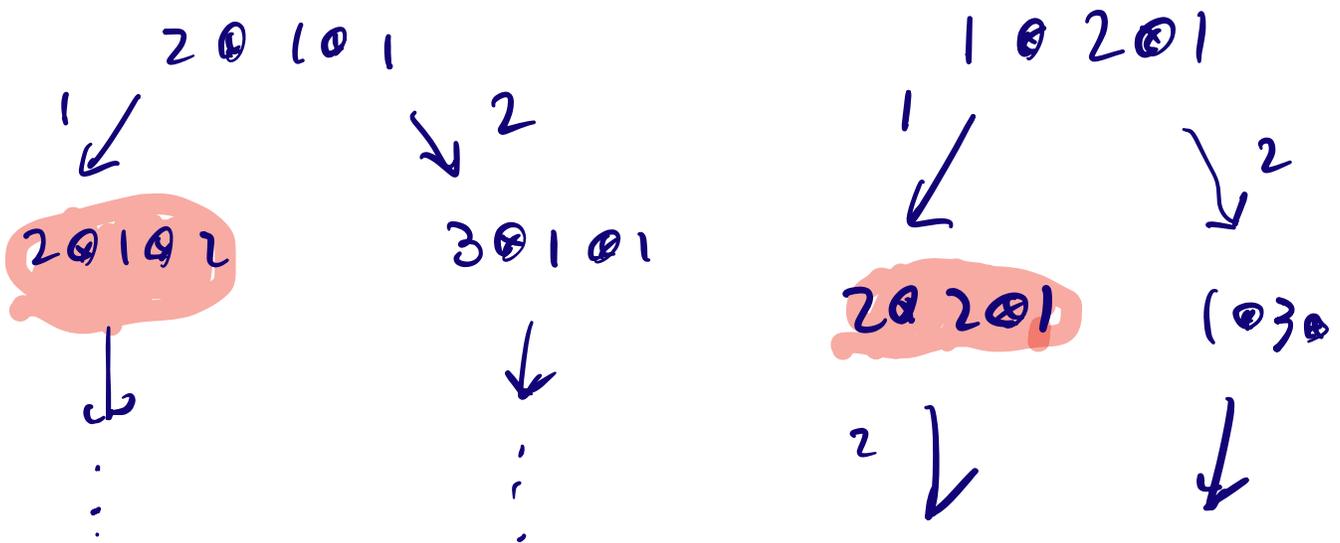
DEF: $T, T' \in \mathbb{Q}^k \mathbb{B}$ ARE

PRACTICALLY EQUIVALENT IF

THEY ARE IN ISOMORPHIC SUBCRYSTALS

$\mathcal{C}, \mathcal{C}' \subset \mathbb{Q}^k \mathbb{B}$ AND THE UNIQUE
 T, T'

ISOMORPHISM $C \rightarrow C'$ TAKES
 $T \rightarrow T'$.



b a c
 $20102 \equiv 20201.$

DONALD KNUTH DETERMINED ALGORITHMICALLY
 WHAT PRACTIC EQUIVALENCE AMOUNTS TO

- (1) If $a < b \leq c$, then bac is Knuth equivalent to bca and vice versa.
- (2) If $a \leq b < c$, then acb is Knuth equivalent to cab and vice versa.